

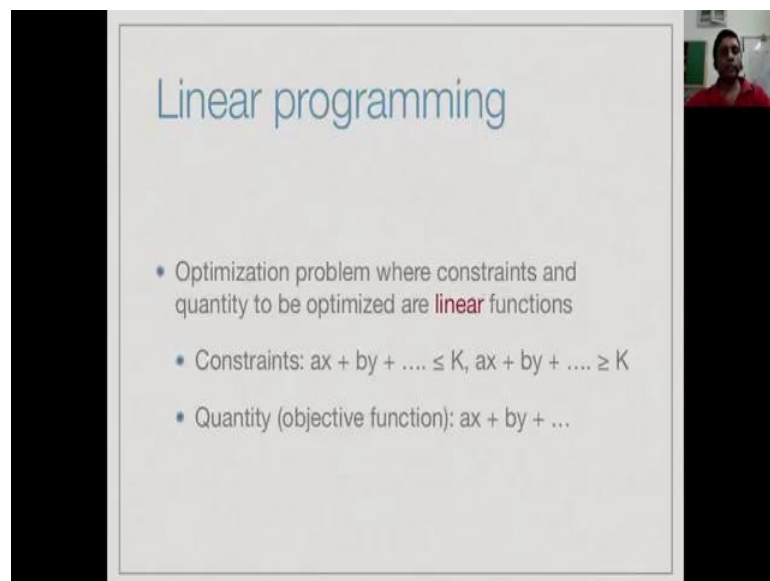
Design and Analysis of Algorithms, Chennai Mathematical Institute
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Department of Computer Science and Engineering,

Week - 08
Module - 02
Lecture – 51

LP Modeling: Production Planning

Let us look at another example of modeling a problem using Linear Programming, again to do with production.

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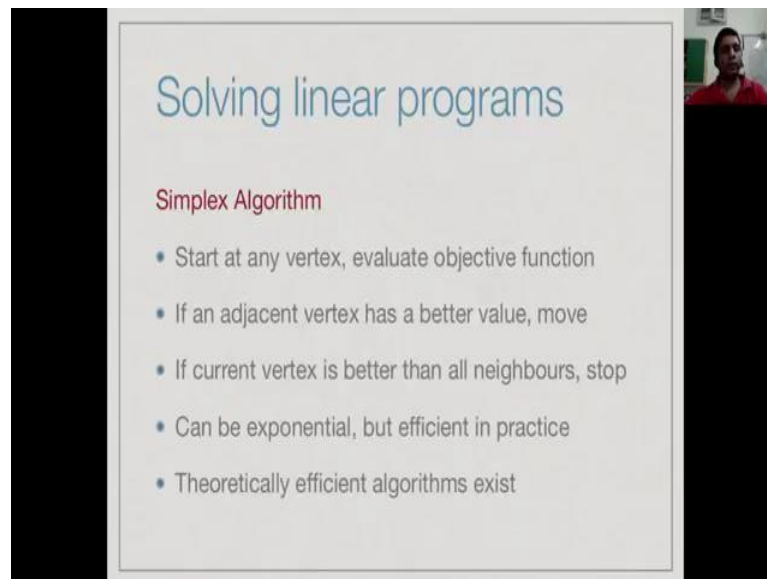


Linear programming

- Optimization problem where constraints and quantity to be optimized are **linear** functions
- Constraints: $ax + by + \dots \leq K$, $ax + by + \dots \geq K$
- Quantity (objective function): $ax + by + \dots$

So, recall that a linear program is an optimization problem, where you have some variables which describe the quantities you want to compute. And you now have linear constraints on these variables, as well as a linear function describing what it is that you want to optimize, maximize or minimize, that is called the objective function.

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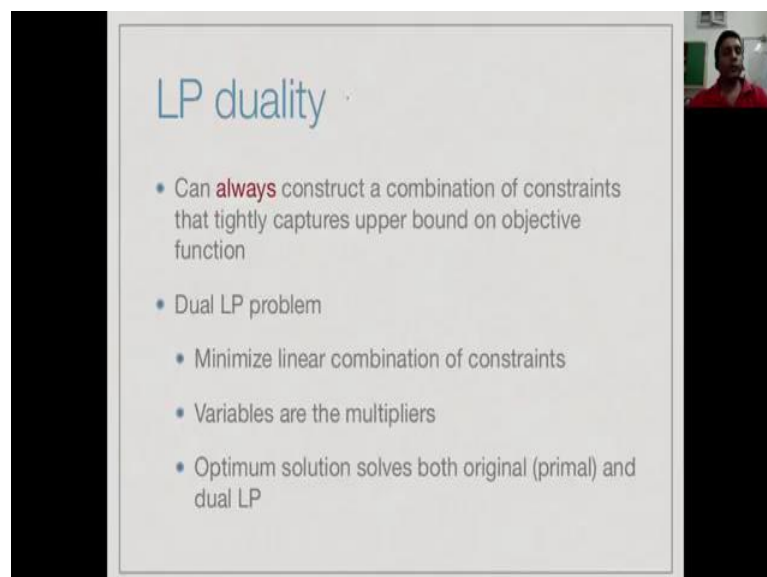
The slide is titled "Solving linear programs" in blue text. Below the title, "Simplex Algorithm" is written in red. A bulleted list follows:

- Start at any vertex, evaluate objective function
- If an adjacent vertex has a better value, move
- If current vertex is better than all neighbours, stop
- Can be exponential, but efficient in practice
- Theoretically efficient algorithms exist

A small video inset in the top right corner shows a person in a red shirt.

And one way to solve a linear program is to think of it geometrically and uses simplex algorithm and the simplex algorithm exploits the fact that the optimum value of a linear program is always at the vertex of the feasible region. So, it starts at some vertex and it keeps going from one vertex to a neighbor, until it finds a vertex whose value is optimum with respect to it is neighbor and the claim is that vertex actually represents the solution.

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The slide is titled "LP duality" in blue text. A bulleted list follows:

- Can **always** construct a combination of constraints that tightly captures upper bound on objective function
- Dual LP problem
 - Minimize linear combination of constraints
 - Variables are the multipliers
- Optimum solution solves both original (primal) and dual LP

A small video inset in the top right corner shows a person in a red shirt.

We also said that we can justify that this is the optimum value by constructing the dual, which has how to combine the constraints together to minimize the combination and the solution which solves both the original and the dual is actually an optimum value.

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Production planning

Handwoven carpets

- 30 employees, each produces 20 carpets a month, salary Rs 20,000 *600 carpets*
- Labour cost is Rs 1000 per carpet
- Monthly demand is seasonal
- Ranges from 440 to 920
- $d_1 \dots d_{12}$ from January to December

So, the next example we are going to look at is the carpet manufacturing company. So, we have a company which makes hand woven carpets and we currently employ 30 employees, each employee produces 20 carpets a month and his pay 20,000 rupees a salary. So, if we just look at the cost per carpet, then we are paying the 1000 rupees to manufacture each carpet.

Now, it turns out that our monthly demand is seasonal, so after looking at some task data we are made an estimate of the amount of carpets that we can expect to sell in each month of the year. So, we have a range of demand from 440 to 920, so remember that you have 30 employees and 20 carpets each, I can make 600 carpets in a given month, this is what I can do with current stake. Now, some times the demand is as low as 440.

So, if I make 600 carpets, now if I have 30 employees and each one is going to produce 20 carpets a month, I am going to get that back. So, I am paying a percent to make carpets it not set to ideal, so every employee will produce 20 carpets a month. Therefore, I will always produce 600 carpets, but I will have 160 pairs which I cannot sell that. But, there are some months where I would have a demand for 320 and now the question of whether I have enough carpets to sell or not.

So, I know the demand from January to December and say these are stored as d_1 for January, d_2 for February and so on to d_{12} for December. So, I have these 12 quantities which are known net marks, these are how many carpets I can expect to sell and each of these marks.

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Coping with varying demand

- Overtime
 - Pay 80% extra, overtime limit 30% per worker
- Hiring and firing
 - Costs Rs 3200 and Rs 4000 per worker
- Store surplus
 - Costs Rs 80 per carpet per month

Handwritten notes in red: Rs 1000/carpet, Rs 1800/carpet overtime rate

Handwritten notes in green: 20/month, ↳ +6

So, now, I have to come up with this strategy to cope with this varying demand, in order to make sure that I lose the least amount of money, because of this coping demand fluctuation. So, one possibility is that whenever I have extra demand I take overtime, so if I pay over time I get workers to work longer hours and make more carpets. But, there are 2 cost to this, the working wage per hour for overtime is typically higher than the regular edge, let us assume it is 80 percent extra.

So, if you remember originally we paid 1000 rupees per carpet as our labor cost, so if I make it in overtime, it is going to cost be instead rupees 1800 per carpet. Because, it is going to take the same amount of time and the person is going to be paid 80 percent more. So, 80 percent of 1800, so it is going to cost 1800 per carpet is the overtime rate, in the other thing is that you cannot of course, expect somebody to work 24 hours a day.

So, there is a limit saying that a worker cannot make more than 30 percent, cannot spend more than 30 percent overtime. So, if they make 20 per month normally, then at most they can go plus 6, a single worker can make at most now including overtime 26 per month. The other option for me is to add or subtract employees, I might want to add new workers, in case I have a higher demand projected a month or I might need to terminate some employees in case my demand probes below 600.

But, these also come to the cost, it goes I have to do some paper work and also I have to give some compensation and so on. So, let us assume that there are some certain cost 3200 is the cost associate in hiring a worker and 4000 rupees is the cost associate in

firing a worker. And finally, I can store carpets when I made exist and sell them later when I have a demand, but storage also cost something. So, let us assume that storing a carpet cost rupees 80 per month.

Because, I have to keep it carefully avoid damage from moisture and so on, these are the various options I can pay overtime. But, up to a limit and it cost more for carpet made by a person who repaid overtime. I can either add or subtract from my work force, but this also comes with the cost each time I do it. And finally, I have a surplus storage cost in case I want to keep my production, hand it over from one month to the next.

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Formulate a linear program

- 1 • w_i : workers in month i , $w_0 = 30$
- 2 • x_i : carpets made in month i
- 3 • o_i : carpets made in overtime in month i
- 4 • h_i : number of workers hired at start of month i
- 5 • f_i : number of workers fired at start of month i
- 6 • s_i : surplus carpets after month i , $s_0 = 0$

72 variables, plus w_0, s_0

Jan Feb ... Dec
1 2 ... 12

So, we want to make a linear program out of this, so we need some variables. So, we are dealing with these 12 months January, February to December. So, it is natural to think of everything as happening within this scope of one month. So, for month i where i ranges from 1 to 12, let us assume that we have W_i workers, initially we have 30 workers. Then, from this 30 we will either hire or fire some to get the workers in the empty, then if hire or fire some to get February and so on.

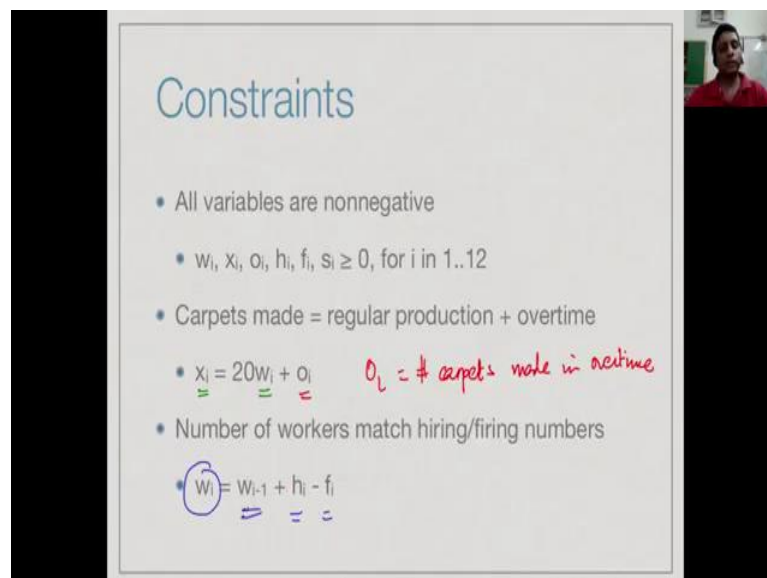
So, each month potentially we have more or less workers, but we start with 30, now these 30 workers will make some number of carpets. So, let us assume that the total number of carpets made in a month is X_i and O_i which is included in X_i in terms of the total. But, O_i is specifically the number of carpets which are made in overtime, because for these there is an extra cost, this is over in other salary we already pay the worker.

And finally, we said that we have this hiring and firing cost, so let h_i be the number of

workers hired at the beginning of a month and f_i be the number of workers fired at the beginning of a month. And let S_i stands for the number of surplus carpets that I have in stock at the end of month i . So, I initially assume that I have an empty variables, so s_0 as 0, but s_1 with say at the end of January how many carpets on my story, s_2 with say at the end of February how many carpets on my story and so on.

So, there are 1, 2, 3, 4, 5, 6 quantities, each of these quantities is there for 12 months. So, as 72 variables plus I have these default initial variables for w_0 or s_0 , so 74 variables. So, this is the lot of things that I have estimate, but it turns out an actually in a linear program this is not a large number, it can be solved quite efficiently by a new of the standard things including simplex.

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Constraints

- All variables are nonnegative
 - $w_i, x_i, o_i, h_i, f_i, s_i \geq 0$, for i in $1..12$
- Carpets made = regular production + overtime
 - $x_i = 20w_i + o_i$ *$O_i = \# \text{ carpets made in overtime}$*
- Number of workers match hiring/firing numbers
 - $w_i = w_{i-1} + h_i - f_i$

So, let us now think of the constraints, so the first constraint to the usual one, which is their every quantity that we are talking about is strictly greater than equal to 0. The second constraints talks about, how the number of carpets made break up into the regular production plus the overtime. So, if I make i carpets in a month, then each worker remember W_i number of workers in my establishment in month i will make 20. So, 20 times W_i what workers in month i produce within the normal time and in addition to that O_i is just the number of carpets made.

So, O_i is not the number of hours overtime, so O_i is the number of carpets. So, each month I make 20 times number of workers plus O_i , now if I start if I had W_i minus 1 workers last month and I hired a few and fired a few, then the total number I have now is

the old number plus the number of hired plus number of hired. So, this is the constraint on how each of the W_i is connected to the previous W_i .

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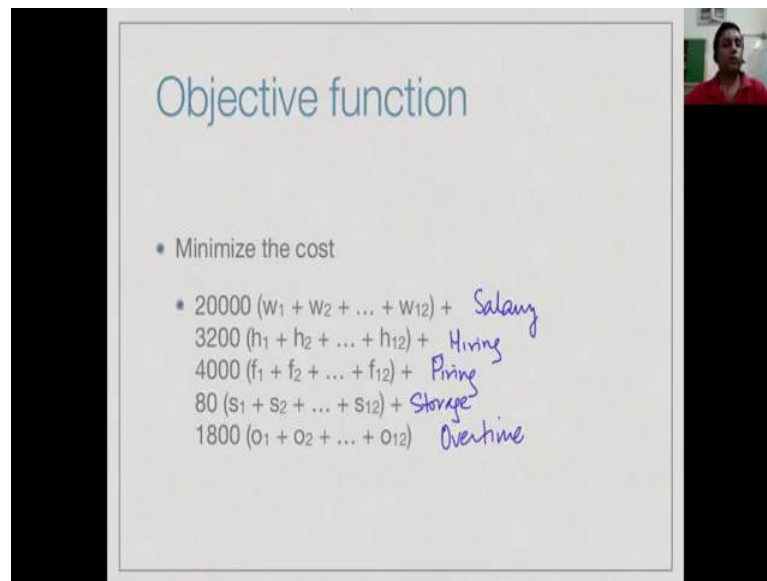
Constraints ...

- Number of stored carpets connected to earlier stock, production, demand
 - $S_i = S_{i-1} + X_i - d_i$ *$d_1 -- d_{12}$ demand*
- Overtime production is at most 6 carpets per worker (30% of regular production)
 - $O_i \leq 6W_i$ *$20 + 30\% = 20 + 6$*

Similarly, the stack is connected to how much I had before I started this month, how much I made this month and how much I sold. So, remember that d_1 to d_{12} was the demand, so this is the fixed quantity, either not variables these are known values. So, in this month I would have sold d_i , but I produce X_i . So, X_i minus d_i is the exits that I produce this month are the deficit as the case may be I started with the surplus S_i minus 1.

So, if I add this month's surplus are defined I get the total surplus at the end of the month. Remember, I cannot have negative carpets, S_i all those things must always be bigger than equal to 0. And finally, as we said before a worker to produces 20 carpets in regular month, if I add now the fact that you can do 30 percent overtime, this is 20 plus 6. So, the number of carpets made in overtime can be at most 6 per worker. So, O_i should be less than or equal to 6 times W_i , so these are all constraints now.

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Objective function

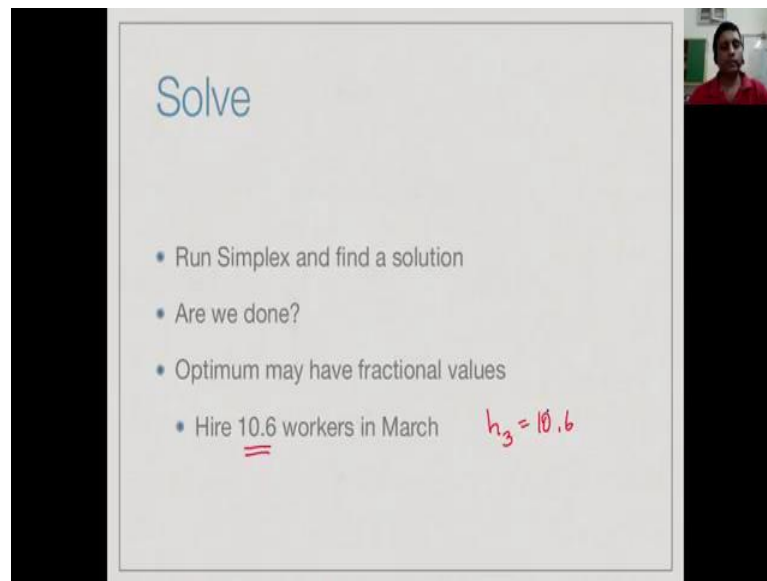
- Minimize the cost
 - $20000 (w_1 + w_2 + \dots + w_{12}) +$ Salary
 - $3200 (h_1 + h_2 + \dots + h_{12}) +$ Hiring
 - $4000 (f_1 + f_2 + \dots + f_{12}) +$ Firing
 - $80 (s_1 + s_2 + \dots + s_{12}) +$ Storage
 - $1800 (o_1 + o_2 + \dots + o_{12}) +$ Overtime

So, now given these constraints what we want to do is minimize the total amount of cost that we are going to put up. So, first this is the regular salary cost, the regular salary cost is 20,000 per worker and we have w_1 workers in month 1, w_2 month 2 and w_{12} month 12. So, we add of the total number of workers, each of the mess way 20,000 rupees. So, this is the total salary before the year, then depending on how many people we hire in each month, for each hire we pay 3000, so this is our hiring pick.

Similarly, for each person that we business we have to pay 4000 first that. So, this is are firing pick, then the number of wrong carpet that have been keep in storage in each month in case 80 rupees per carpet cost. So, have to multiply the total storage by 80 and finally, when I pay overtime and paying our number the salary. So, each carpet made in overtime, remember cost 1800 rupees as suppose to 1000 rupees which is already accounted for in salary, so this is my overtime cost.

So, this says that we can take this fairly complicated looking question about production with the demand and hiring and inspiring and productivity all that and set it out with some 72 plus 274 variables and complex cost function and feed it to simplex and get an answer.

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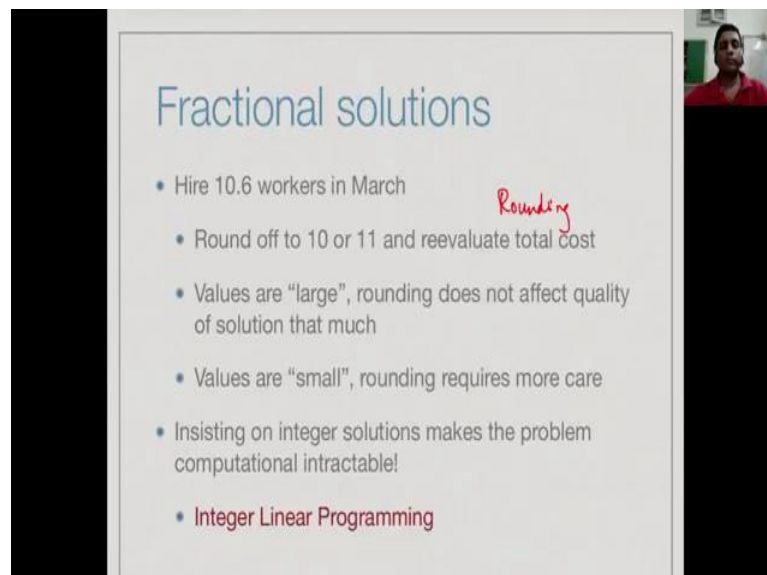
The slide is titled "Solve" in blue. It contains a bulleted list of steps:

- Run Simplex and find a solution
- Are we done?
- Optimum may have fractional values
- Hire 10.6 workers in March

Handwritten in red next to the last bullet is $h_3 = 10.6$.

So, we done simplex and we find the solution, so are we done, so turns out that we might with an answer which is not something that we can actually used. For instance, we might it an answer will says h_3 is equal to 10.6. So, this is the, we must hire a fractional number of people in a month. Now, obviously, we cannot do this at the same time we have not any of our constraints express that the values must be integers.

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The slide is titled "Fractional solutions" in blue. It contains a bulleted list of points:

- Hire 10.6 workers in March
- Round off to 10 or 11 and reevaluate total cost
- Values are "large", rounding does not affect quality of solution that much
- Values are "small", rounding requires more care
- Insisting on integer solutions makes the problem computational intractable!
- Integer Linear Programming

Handwritten in red next to the first bullet is the word "Rounding".

So, what can be do about this, well we can take this 10.6 and we can look for the errors integer 10 or a 11 and reevaluate what happens to the cost, if we make it even integer. So, this is called integer rounding, now if the values are big now direct 10, 11 the two digit value say, then rounding is relatively small displacement and the quality of the optimum

will not change must based on it is. But, if the numbers are small, if I am rounding between say I am going to take it from 0.51 hour 0 of 1.52 or 1 then the effect of the rounding is quite can be quite task.

So, this is the general problem with linear programming which is linear programming cannot guarantee that you have the integer solutions, when you want an integer solution you can use rounding to achieve an integer solution, but you have to be careful that the rounding actually gives you one optimum. Now, why not just insist that you want an integer solution. So, can you not set up this same problem and solve it, but require the solutions to be integers, unfortunately this terms out to be quite a hard probe.

So, we have found we have claim rather that when we set up general linear programming problem, we can solve them efficiently. So, simplex is not an efficient solution, but an effective solution, but there are interior point methods and other polynomial time algorithms for linear programming. But, if you change the rules of the game and say I do not want to arbitrary solution; that means, constraints, but I want to one we are all the values that I get an integers, then we have the, so call integer linear programming.

And integer linear programming unfortunately is not solvable efficiently or it is not known divisible efficiently. So, the best we can hope to do for integer linear programming is to actually try it turn into an arbitrary linear program and some of interpret the answers or integers by doing appropriate rounding.